

Games solved: Now and in the future

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Abstract

- Which game characters are predominant when the solution of a game is the main target?
 - It is concluded that **decision complexity** is more important than **state-space complexity**.
 - There is a trade-off between **knowledge-based methods** and **brute-force methods**.
 - There is a clear correlation between the first-player's **initiative** and the necessary effort to solve a game.

Definitions I

- **Domain: two-person zero-sum games with perfect information.**
 - Zero-sum means one player's loss is exactly the other player's gain, and vice versa. There is no way for both players to win at the same time.
- **Game-theoretic value** of a game: the outcome, i.e., win, loss or draw, when all participants play optimally.
- **Classification of games' solutions according to L.V. Allis in 1994.**
 - **Ultra-weakly solved:** the game-theoretic value of the initial position has been determined.
 - **Weakly solved:** for the initial position a strategy has been determined to achieve the game-theoretic value against any opponent.
 - ▷ *The strategy must be efficient and practical in terms of resource usage.*
 - **Strongly solved:** a strategy has been determined for all legal positions.

Definitions II

- **State-space** complexity of a game: the number of all legal positions in a game.
- **Game-tree** (or **decision**) complexity of a game: the number of all leaf nodes in a *solution search tree*.
 - ▷ *A solution search tree is a tree where the game-theoretic value of the root position can be decided.*
- A **convergence** game: the size of the state space decreases as the game progress.
 - Start with many pieces on the board and pieces are gradually removed during the course of the game.
- A **divergence** game: the size of the state space increases as the game progress.
 - Start with an empty board and pieces are gradually added during the course of the game.
- **Initiative**: the right to move first
- A **fair** game: the game-theoretic value is draw and both players have roughly an equal probability on making a mistake.

Prediction made in 1990

- Predictions were made at 1990 for the year 2000 concerning the expected playing strength of computer programs.

| solved | over champion | world champion | grand master | amateur |
|-------------------|----------------|------------------|---------------|------------|
| Connect-four | Checkers (8*8) | Chess | Go (9*9) | Go (19*19) |
| Qubic | Renju | Draughts (10*10) | Chinese chess | |
| Nine Men's Morris | Othello | | Bridge | |
| Go-Moku | Scrabble | | | |
| Awari | Backgammon | | | |

- ▷ *Over champion means definitely over the best human player.*
- ▷ *World champion means equaling to the best human player.*
- ▷ *Grand master means equaling to most human players.*

Game space

- A double dichotomy of the game space

↑
 $\log \log(\text{state-space complexity})$

| | |
|---|---|
| category 3 if solvable at all, then by knowledge-based methods | category 4 unsolvable by any method |
| category 1 solvable by any method | category 2 if solvable at all, then by brute-force methods |

$\log \log(\text{game-tree complexity}) \rightarrow$

Questions to be researched

- Can perfect knowledge obtained from solved games be translated into rules and strategies which human beings can assimilate?
- Are such rules generic, or do they constitute a multitude of ad hoc recipes?
- Can methods be transferred between games?
 - More specifically, are there generic methods for all category- n games, or is each game in a specific category a law unto itself?

Convergent games

- Can be solved by the method of **endgame databases** if we can enumerate and store all possible positions.
- Problems solved:
 - Nine Men's Morris: in the year 1995, a total of 7,673,759,269 states.
 - ▷ *The game theoretic value is draw.*
 - Mancala games
 - ▷ *Awari: in the year 2002.*
 - ▷ *Kalah: in the year 2000 upto, but not equal, Kalah(6,6)*
 - Checkers: in the year 1994
 - ▷ *By combining Endgame databases, middle-game databases and verification of opening analysis.*
 - ▷ *Solved the so called 100-year position.*
 - Chess endgames
 - Chinese chess endgames

Divergent games

- **Connection games**
 - Connect-four ($6*7$)
 - Qubic ($4*4*4$)
 - Go-Moku ($15*15$)
 - Renju
 - k -in-a-row games
 - Hex ($10*10$ or $11*11$)
- **Polynmino games**
 - Pentominoes
 - Domineering
- **Othello**
- **Chess**
- **Chinese chess**
- **Shogi**
- **Go**

Connection games I

■ Connect-four (6*7)

- Solved by J. Allen in 1989 using a brute-force depth first search with alpha-beta pruning, a transposition table, and killer-move heuristic.
- Also solved by L.V. Allis in 1988 using a knowledge-based approach by combining 9 strategic rules that identified potential **threats** of the opponent.
 - ▷ *Threats are something like forced moves or moves you have little choices.*
 - ▷ *Moves with predictable counter-moves.*
- It is a first-player win.
- Weakly solved on a SUN 4 workstation using 300+ hours.

■ Qubic (4*4*4)

- A three-dimensional version of Tic-Tac-Toe.
- Solved in 1980 by O. Patashnik by combining the usual depth-first search with expert knowledge for ordering the moves.

Connection games II

- **Go-Moku (15*15)**
 - First-player win
 - Weakly solved by L.V. Allis in 1995 using a combination of threat-space search and database construction.
- **Renju**
 - Does not allow the first player to play certain moves.
 - An *asymmetric* game.
 - Weakly solved by Wágner and Virág in 2000 by combining search and knowledge.
 - ▷ *It is still first-player win.*
 - Took advantage of an iterative-deepening search based on threat sequences up to 17 plies.
- **k -in-a-row games**
 - mnk -Game: a game playing on a board of m rows and n columns with the goal of obtaining a straight line of length k .
 - Variations: first ply picks only one stone, the rest picks two stones in a ply.
 - ▷ *Try to balance the advantage of the initiative!*

Hex ($10*10$ or $11*11$)

■ Properties:

- It is a finite game.
- It is not possible for both players to win at the same.
- Exactly one of the players can win.
 - ▷ *A 2-D induction and topological argument.*
 - ▷ *True for Hex $2*2$*
 - ▷ *Assume it is true for any Hex $i * j$, where $i \leq n$ or $j \leq m$.*
 - ▷ *Induction step: try to prove this is true on Hex $n * m$*
 - ▷ *Delete the first row or the last row give you two white chains w_1 and w_2 , respectively.*
 - ▷ *Delete the first column or the last column give you two black chains b_1 and b_2 , respectively.*
 - ▷ *Delete first row, last row, first column and last column gives you either a white chain w_3 or a black chain b_3 .*
 - ▷ *Either w_3 intersects with b_1 or b_2 , or b_3 intersects with w_1 or w_2 ; both are contradicting statements.*
 - ▷ *Other arguments in graph theory exist.*

Strategy-stealing argument

- The *unrestricted form is a first-player win*.
 - Using the “strategy-stealing” argument made by John Nash in 1949.
 - ▷ *If there is a winning strategy for the second player, the first player can still win by making an arbitrary first move and using the second-player strategy from then on.*
 - ▷ *If using the second-player strategy requires playing the chosen first move or any move played before, then make another arbitrary move.*
 - ▷ *An arbitrary extra move can never be a disadvantage in Hex.*
 - ▷ *This is not true for every games.*
 - ▷ *Not a constructive proof.*
 - This argument works for any symmetry games when an arbitrary extra move can never be a disadvantage.
 - The **one-move-equalization** rule: one player plays an opening move and the other player then has to decide which color to play for the remainder of the game.
 - The revised version is a second-player win (ultra-weakly).

Solutions to HEX

- Hex exhibits considerable mathematical structure.
- Hex has been proved to be PSPACE-complete by Even and Tarjan in 1976 by converting it to a Shannon switching game.
- The state-space and decision complexities are comparable to those of Go on equally-sized boards.
- (Weakly or strongly) solved on $6*6$ boards in 1994.
- Maybe possible to solve the $7*7$ case.
- Not likely to solve the $8*8$ version without fundamental breakthroughs.

More divergent games

- **Polynmino games: placing 2-D pieces that forms a connected subset of a square grid on a grid.**
 - Pentominoes
 - Domineering
 - Games on smaller boards are solved.
- **Othello**
 - M. Buro's LOGISTELLO beat the resigning World Champion by 6-0 in 1997.
 - Weakly solved on 6*6 boards by J. Feinstein in 1993.
- **Chess**
 - DEEP BLUE beat the human World Champion in 1997.
- **Chinese chess**
 - Still in progress,
 - Professional 7-dan in 2007.
- **Shogi**
 - Still in progress,
 - Professional 2-dan in 2007.
- **Go**
 - Still in progress,
 - Amateur 4 kyu.

State-space versus game-tree size

- In 1994, the boundary of solvability by complete enumeration was set at 10^{11} .
 - The current estimation is about 10^{13} .

Table of complexity

| Game | $\log_{10}(\text{state-space})$ | $\log_{10}(\text{game-tree size})$ |
|-------------------|---------------------------------|------------------------------------|
| Nine Men's Morris | 10 | 50 |
| Pentominos | 12 | 18 |
| Awari | 12 | 32 |
| Kalak(6,4) | 13 | 18 |
| Connect-four | 14 | 21 |
| Domineering (8*8) | 15 | 27 |
| Dakon-6 | 15 | 33 |
| Checkers | 21 | 31 |
| Othello | 28 | 58 |
| Qubic | 30 | 34 |
| Draughts | 30 | 54 |
| Chess | 46 | 123 |
| Chinese chess | 48 | 150 |
| Hex (11*11) | 57 | 98 |
| Shogi | 71 | 226 |
| Renju (15*15) | 105 | 70 |
| Go-Moku (15*15) | 105 | 70 |
| Go (19*19) | 172 | 360 |

Brute-force versus knowledge-based methods

- Games with both a relative low state-space complexity and a low game-tree complexity have been solved by both methods.
 - Category 1
 - Connect-four and Qubic
- Games with a relative low state-space complexity have mainly been solved with brute-force methods
 - Category 2
 - Namely by constructing endgame database
 - Nine Men's Morris
- Games with a relative low game-tree-complexities have mainly been solved with knowledge-based methods.
 - Category 3
 - Namely, by intelligent (heuristic) searching
 - Sometimes, with the helps of endgame databases
 - Go-Moku, Renju, and k -in-a-row games

The advantage of the initiative

- **Theorem (or arguments) made by Singmaster in 1981: The first-player has advantages.**
 - **Two kinds of positions**
 - ▷ *P-positions: the previous player can force a win.*
 - ▷ *N-positions: the next player can force a win.*
 - **Arguments**
 - ▷ *For the first player to have a forced win, just one of the moves must lead to a P-position.*
 - ▷ *For the second player to have a forced win, all of the moves must lead to N-positions.*
 - ▷ *It is easier to the first player to have a forced win assuming all positions are randomly distributed.*
 - ▷ *Can be easily extended to games with draws.*
 - **Remarks:**
 - ▷ *On small boards, the second player is able to draw or even to win for certain games.*
 - ▷ *Cannot be applied to the infinite board.*

The advantage of the initiative in practice

- A connection $mn1$ -game.
- Mentioned in S.K Chang's famous 1978 novel: Chi Wang.
- Need to offset the initiative
 - Enforce rules so that the first player can win by selective patterns.
 - Swapping: a player makes the first move, the second player decides the color to play thereafter.
 - The first move plays one stone, the rest plays two stones.
 - The first player uses less time.
 - Must be simple and try to be as fair as possible.

How to make use of initiative

- **A potential universal strategy for winning a game:**
 - Try to obtain a small initiative
 - ▷ *The opponent must react adequately on the moves played by the player.*
 - To reinforce the initiative the player searches for threats, and even a sequence of threats using an evaluation function E .
- **Threat-space search**
 - Search for threats only!

Methods developed for solving games

■ Brute-force methods

- Retrograde analysis
- Enhanced transposition-table methods

■ Knowledge-based methods

- Threat-space search and λ -search
- Proof-number search
- Depth-first proof-number search
- Pattern search

- ▷ *To search for **threat patterns**, which are collections of cells in a position P .*
- ▷ *A threat pattern can be thought of as representing the **relevant area** on the board, an area that human players commonly identify when analyzing a position.*

Conclusions

- **The knowledge-based methods mostly inform us on the structure of the game, while exhaustive enumeration rarely does.**
- **Many ad-hoc recipes are produced currently.**
 - The database can be used as a corrector of strategies formulated by human experts.
- **It may be hopeful to use data mining techniques to obtain cross-game methods.**
 - Currently not very successful.

New predictions

- Predictions were made for the year 2010 concerning the expected playing strength of computer programs.

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|----------------|------------------|----------------|--------------|------------|
| Awari | Chess | Go (9*9) | Bridge | Go (19*19) |
| Othello | Draughts (10*10) | Chinese chess | Shogi | |
| Checkers (8*8) | Scrabble | Hex | | |
| | Backgammon | Amazons | | |
| | Lines of Action | | | |

References and further readings

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